# 01204211 Discrete Mathematics <br> Lecture 1b: Implications and equivalences 

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## This lecture covers:

- More connectives: implications and equivalences


## Review (1)

- A proposition is a statement which is either true or false.
- We can use variables to stand for propositions, e.g., $P=$ "today is Tuesday".
- We can use connectives to combine variables to get propositional forms.
- Conjunction: $P \wedge Q$ (" $P$ and $Q$ "),
- Disjunction: $P \vee Q$ (" $P$ or $Q$ "), and
- Negation: $\neg P($ ( $n$ ot $P$ " $)$


## Review (2)

To represents values of propositional forms, we usually use truth tables.

## And/Or/Not

| $P$ | $Q$ | $P \wedge Q$ | $P \vee Q$ | $\neg P$ |
| :---: | :---: | :---: | :---: | :---: |
| $T$ | $T$ | $T$ | $T$ | $F$ |
| $T$ | $F$ | $F$ | $T$ |  |
| $F$ | $T$ | $F$ | $T$ | $T$ |
| $F$ | $F$ | $F$ | $F$ |  |

## Quick check 1

As we said before, the truth value of propositional forms may not depend on the values of its variables. As you can see in this exercise.
Use a truth table to find the values of (1) $P \wedge \neg P$ and (2) $P \vee \neg P$.

## And/Or/Not

| $P$ | $\neg P$ | $P \wedge \neg P$ | $P \vee \neg P$ |
| :---: | :---: | :---: | :---: |
| $T$ | $F$ | $F$ | $T$ |
| $F$ | $T$ | $F$ | $T$ |

Note that $P \wedge \neg P$ is always false and $P \vee \neg P$ is always true. A propositional form which is always true regardless of the truth values of its variables is called a tautology. On the other hand, a propositional form which is always false regardless of the truth values of its variables is called a contradiction.

## Implications ${ }^{1}$

Given $P$ and $Q$, an implication

$$
P \Rightarrow Q
$$

stands for "if $P$, then $Q$ ". This is a very important propositional form.
It states that "when $P$ is true, $Q$ must be true". Let's try to fill in its truth table:

Implications

| $P$ | $Q$ | $P \Rightarrow Q$ |
| :---: | :---: | :---: |
| $T$ | $T$ | $T$ |
| $T$ | $F$ | $F$ |
| $F$ | $T$ | $T$ |
| $F$ | $F$ | $T$ |

[^0]
## What?

- Yes, when $P$ is false, $P \Rightarrow Q$ is always true no matter what truth value of $Q$ is.
- We say that in this case, the statement $P \Rightarrow Q$ is vacuously true.
- You might feel a bit uncomfortable about this, because in most natural languages, when we say that if $P$, then $Q$ we sometimes mean something more than that in the logical expression " $P \Rightarrow Q$."


## One explanation

- But let's look closely at what it means when we say that: if $P$ is true, $Q$ must be true.
- Note that this statement does not say anything about the case when $P$ is false, i.e., it only considers the case when $P$ is true.
- Therefore, having that $P \Rightarrow Q$ is true is OK with the case that (1) $Q$ is false when $P$ is false, and (2) $Q$ is true when $P$ is false.
- This is an example when mathematical language is "stricter" than natural language.


## Noticing if-then

We can write "if $P$, then $Q$ " for $P \Rightarrow Q$, but there are other ways to say this. E.g., we can write (1) $Q$ if $P$, (2) $P$ only if $Q$, or (3) when $P$, then $Q$.

## Quick check 2

For each of these statements, define propositional variables representing each proposition inside the statement and write the proposition form of the statement.

- If you do not have enough sleep, you will feel dizzy during class.
- If you eat a lot and you do not have enough exercise, you will get fat.
- You can get A from this course, only if you work fairly hard.


## Only-if

Let $P$ be "you get A from this course."
Let $Q$ be "you work fairly hard."
Let $R$ be "You can get A from this course, only if you work fairly hard."
Let's think about the truth values of $R$.
Only if you work fairly hard.

| $P$ | $Q$ | $R$ |
| :--- | :--- | :--- |
| $T$ | $T$ |  |
| $T$ | $F$ |  |
| $F$ | $T$ |  |
| $F$ | $F$ |  |

Thus, $R$ should be logically equivalent to $P \Rightarrow Q$. (We write $R \equiv P \Rightarrow Q$ in this case.)

## If and only if: $(\Leftrightarrow)$

Given $P$ and $Q$, we denote by

$$
P \Leftrightarrow Q
$$

the statement " $P$ if and only if $Q$." It is logically equivalent to

$$
(P \Leftarrow Q) \wedge(P \Rightarrow Q)
$$

i.e., $P \Leftrightarrow Q \equiv(P \Leftarrow Q) \wedge(P \Rightarrow Q)$.

Let's fill in its truth table.

| $P$ | $Q$ | $P \Rightarrow Q$ | $P \Leftarrow Q$ | $P \Leftrightarrow Q$ |
| :---: | :---: | :---: | :---: | :---: |
| $T$ | $T$ |  |  |  |
| $T$ | $F$ |  |  |  |
| $F$ | $T$ |  |  |  |
| $F$ | $F$ |  |  |  |

## An implication and its friends

When you have two propositions

- $P=$ "I own a cell phone", and
- $Q=$ "I bring a cell phone to class".

We have

- an implication $P \Rightarrow Q \equiv$
"If I own a cell phone, I'll bring it to class",
- its converse $Q \Rightarrow P \equiv$
"If I bring a cell phone to class, I own it", and
- its contrapositive $\neg Q \Rightarrow \neg P \equiv$
"If I do not bring a cell phone to class, I do not own one".


## Quick check 3

Let's consider the following truth table:

| $P$ | $Q$ | $P \Rightarrow Q$ | $Q \Rightarrow P$ | $\neg Q \Rightarrow \neg P$ |
| :---: | :---: | :---: | :---: | :---: |
| $T$ | $T$ |  |  |  |
| $T$ | $F$ |  |  |  |
| $F$ | $T$ |  |  |  |
| $F$ | $F$ |  |  |  |

Do you notice any equivalence?
Right, $P \Rightarrow Q \equiv \neg Q \Rightarrow \neg P$.


[^0]:    ${ }^{1}$ Materials in this lecture are mostly from Berkeley CS70's lecture notes.

